Artificial UCLouvain Intelligence & Algorithms **Anytime Weighted Model Counting with Approximation Guarantees For Probabilistic** Inference

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#### **Probabilistic Inference: Bayesian Networks**

press, 2009.



#### **Probabilistic Inference: Bayesian Networks**



### Probabilistic Inference: Query on Bayesian Networks



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Chavira, Mark, and Adnan Darwiche. "On probabilistic inference by weighted model counting." Artificial Intelligence 172.6-7 (2008): 772-799.



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Chavira, Mark, and Adnan Darwiche. "On probabilistic inference by weighted model counting." *Artificial Intelligence* 172.6-7 (2008): 772-799. Valiant, Leslie G. "The complexity of enumeration and reliability problems." *siam Journal on Computing* 8.3 (1979): 410-421.





## Why Model Counting? "One solver to solve them all"



# Why Model Counting? "One solver to solve them all"





## Probabilistic Inference for Large language Models



Computing exact probabilities "That's a no-no" - Donald Trump

Computing **approximate** probabilities "Yes we can" - Barack Obama

 $\tilde{p}$ 



Approximated probability

$$rac{p^{\star}}{1+arepsilon}\leq ext{ } ilde{p}$$

True probability





Epsilon guarantee bounds the true probability with relative error

$$rac{p^{\star}}{1+arepsilon} \leq ~~ ilde{p}~\leq (1+arepsilon)p^{\star}$$

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$$P\left[rac{p^{\star}}{1+arepsilon} \leq ilde{p} \leq (1+arepsilon)p^{\star}
ight] \geq 1-\delta$$

Epsilon-delta guarantee bounds the true count probabilistically

Classical DFS Algorithm









	Anytime	Guarantees	Lower bound	Upper bound
d4/gpmc/ exactMC/	X	Exact on completion	X	X

Lagniez, Jean-Marie, and Pierre Marquis. "A recursive algorithm for projected model counting." *Proceedings of the AAAI Conference on Artificial Intelligence*. Vol. 33. No. 01. 2019.

Lagniez, Jean-Marie, and Pierre Marquis. "An Improved Decision-DNNF Compiler." IJCAI. Vol. 17. 2017.

Ryosuke Suzuki, Kenji Hashimoto, and Masahiko Sakai. Improvement of projected model counting solver with component decomposition using SAT solving in components.

Yong Lai, Kuldeep S. Meel, and Roland HC Yap. The power of literal equivalence in model counting. In Proceedings of the AAAI Conference on Artificial Intelligence, volume 35, 2021.

	Anytime	Guarantees	Lower bound	Upper bound
d4/gpmc/ exactMC/…	X	Exact on completion	X	X
ApproxMC/ WeightMC		$(arepsilon,\delta)$		

Supratik Chakraborty, Kuldeep S. Meel, and Moshe Y. Vardi. Algorithmic improvements in approximate counting for probabilistic inference: From linear to logarithmic SAT calls. In IJCAI, 2016.

Mate Soos, Stephan Gocht, and Kuldeep S. Meel. Tinted, Detached, and Lazy CNF-XOR solving and its Applications to Counting and Sampling. In CCAV, 2020.

Mate Soos and Kuldeep S. Meel. BIRD: Engineering an efficient CNF-XOR SAT solver and its applications to approximate model counting. In Proceedings of the AAAI Conference on Artificial Intelligence

Supratik Chakraborty, Daniel Fremont, Kuldeep Meel, Sanjit Seshia, and Moshe Vardi. Distribution-aware sampling and weighted model counting for SAT. In AAAI, 2014.

	Anytime	Guarantees	Lower bound	Upper bound
d4/gpmc/ exactMC/…	X	Exact on completion	X	X
ApproxMC/ WeightMC		$(arepsilon,\delta)$		
PartialKC / SampleSAT		Probability on the LB		

Yong Lai, Kuldeep S. Meel, and Roland HC Yap. Fast Converging Anytime Model Counting. In Proceedings of the AAAI Conference on Artificial Intelligence, volume 37, 2023.

	Anytime	Guarantees	Lower bound	Upper bound
d4/gpmc/ exactMC/…	X	Exact on completion	X	X
ApproxMC/ WeightMC		$(arepsilon,\delta)$		
PartialKC / SampleSAT		Probability on the LB		
Toulbar		(arepsilon,0)		

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ApproxMC/ WeightMC		$(arepsilon,\delta)$		
PartialKC / SampleSAT		Probability on the LB		
Toulbar		(arepsilon,0)		
Our contribution		(arepsilon,0)		

	Anytime	Guarantees	Lower bound	Upper bound
d4/gpmc/ exactMC/…	X	Exact on completion	X	X
ApproxMC/ WeightMC		$(arepsilon,\delta)$		
PartialKC / SampleSAT		Probability on the LB		
Toulbar		(arepsilon,0)		
Schlandals		(arepsilon,0)		

# **DFS** with Bounds

Lower bound on the (sub-)problem(sub-)problem

## DFS with Bounds

$$D_3 = \{ rac{d_1}{d_3}, d_3^2, d_3^3, d_3^4 \}$$
  $[0.0, 1.0]$   
 $D_3 = \{ rac{d_1}{d_3}, d_3^2, d_3^3, d_3^4 \}$   
 $D_2$   
 $D_3 \neq d_1^1$   
 $D_3 \neq d_3^1$   
 $P(D_3 = d_3^1) = 0.2$ 

### **DFS** with Bounds












# Fifty shades of Approximate Inference

	Anytime	Guarantees	Lower bound	Upper bound
d4/gpmc/ exactMC/…	X	Exact on completion	X	X
ApproxMC/ WeightMC		$(arepsilon,\delta)$		
PartialKC / SampleSAT		Probability on the LB		
Toulbar		(arepsilon,0)		
Schlandals		(arepsilon,0)		

$$rac{p^{\star}}{1+arepsilon} \leq ilde{p} \leq (1+arepsilon)p^{\star}$$





$$egin{aligned} &rac{p^{\star}}{1+arepsilon} & \leq ilde{p} \leq (1+arepsilon) p^{\star} \ & \downarrow ? \ & f(lb,ub) & ub \leq lb(1+arepsilon)^2 \Rightarrow rac{p^{\star}}{1+arepsilon} \leq \sqrt{lb imes ub} \leq (1+arepsilon) p^{\star} \end{aligned}$$

$$\begin{array}{c|c} p^{\star} & \leq \tilde{p} \leq (1+\varepsilon)p^{\star} \\ & & \\ & & \\ & & \\ \end{array} \begin{array}{c} ? \\ f(lb,ub) \end{array} \begin{array}{c} \text{Proof in the paper !} \\ ub \leq lb(1+\varepsilon)^2 \Rightarrow \frac{p^{\star}}{1+\varepsilon} \leq \sqrt{lb \times ub} \leq (1+\varepsilon)p^{\star} \end{array}$$

$$egin{aligned} & p^{\star} & \leq ilde{p} \leq (1+arepsilon) p^{\star} \ & \downarrow ? \ & f(lb,ub) & ub \leq lb(1+arepsilon)^2 \Rightarrow rac{p^{\star}}{1+arepsilon} \leq \sqrt{lb imes ub} \leq (1+arepsilon) p^{\star} \end{aligned}$$

$$arepsilon_{lb,ub} = \sqrt{rac{ub}{lb}} - 1$$

Lower- and upper-bounds induced a minimal approximation factor !

# Fifty shades of Approximate Inference

	Anytime	Guarantees	Lower bound	Upper bound
d4/gpmc/ exactMC/…	X	Exact on completion	X	X
ApproxMC/ WeightMC		$(arepsilon,\delta)$		
PartialKC / SampleSAT		Probability on the LB		
Toulbar		(arepsilon,0)		
Schlandals		(arepsilon,0)		

	P(query)	P(que
$w_{I_1} = 0.0035$	0.0035	
$w_{I_2}=0.06$		
$w_{I_3}=0.15$		
$w_{I_4}=0.2$		
$w_{I_5}=0.07$		
$w_{I_6} = 0.1165$		
$w_{I_7}=0.4$		

 $P(query) = 1 - P(\neg query)$ 

1.0



	P(query)	$P(query) = 1 - P(\neg query)$
$w_{I_1} = 0.0035$		
$w_{I_2}=0.06$		
$w_{I_3}=0.15$	0.0635	1 - 0.15 = 0.85
$w_{I_4}=0.2$		
$w_{I_5} = 0.07$		
$w_{I_6} = 0.1165$		
$w_{I_7}=0.4$		

	P(query)	$P(query) = 1 - P(\neg q$
$w_{I_1} = 0.0035$		
$w_{I_2}=0.06$		
$w_{I_3}=0.15$		
$w_{I_4}=0.2$	0.2635	1 - 0.15 = 0.85
$w_{I_5}=0.07$		
$w_{I_6} = 0.1165$		
$w_{I_7}=0.4$		

 $\neg query)$ 

 $w_{I_1} = 0.0035$  $w_{I_2} = 0.06$  $w_{I_3} = 0.15$  $w_{I_4}=0.2$  $w_{I_5} = 0.07$  $w_{I_6} = 0.1165$  $w_{I_7}=0.4$ 

P(query)  $P(query) = 1 - P(\neg query)$ 

	P(query)	$P(query) = 1 - P(\neg query)$		P(query)	$P(query) = 1 - P(\neg query)$
$w_{I_1} = 0.0035$	0.0035	1.0	$w_{I_7}=0.4$	0.0	0.6
$w_{I_2} = 0.06$			$w_{I_4}=0.2$		
$w_{I_3} = 0.15$			$w_{I_3} = 0.15$		
$w_{I_4}=0.2$			$w_{I_6} = 0.1165$		
$w_{I_5} = 0.07$			$w_{I_5} = 0.07$		
$w_{I_6} = 0.1165$			$w_{I_2} = 0.06$		
$w_{I_7}=0.4$			$w_{I_1} = 0.0035$		

	P(query)	$P(query) = 1 - P(\neg query)$		P(query)	$P(query) = 1 - P(\neg query)$
$w_{I_1} = 0.0035$			$w_{I_7}=0.4$		
$w_{I_2} = 0.06$	0.0635	1.0	$w_{I_4}=0.2$	0.2	0.6
$w_{I_3}=0.15$			$w_{I_3} = 0.15$		
$w_{I_4}=0.2$			$w_{I_6} = 0.1165$		
$w_{I_5}=0.07$			$w_{I_5} = 0.07$		
$w_{I_6} = 0.1165$			$w_{I_2}=0.06$		
$w_{I_7}=0.4$			$w_{I_1} = 0.0035$		
	$w_{I_1} = 0.0035$ $w_{I_2} = 0.06$ $w_{I_3} = 0.15$ $w_{I_4} = 0.2$ $w_{I_5} = 0.07$ $w_{I_6} = 0.1165$ $w_{I_7} = 0.4$	$P(query)$ $w_{I_1} = 0.0035$ $w_{I_2} = 0.06$ $w_{I_3} = 0.15$ $w_{I_4} = 0.2$ $w_{I_5} = 0.07$ $w_{I_6} = 0.1165$ $w_{I_7} = 0.4$	$P(query) \qquad P(query) = 1 - P(\neg query)$ $w_{I_1} = 0.0035$ $w_{I_2} = 0.06$ $w_{I_3} = 0.15$ $w_{I_4} = 0.2$ $w_{I_5} = 0.07$ $w_{I_6} = 0.1165$ $w_{I_7} = 0.4$	$P(query)  P(query) = 1 - P(\neg query)$ $w_{I_1} = 0.0035$ $w_{I_2} = 0.06$ $w_{I_3} = 0.15$ $w_{I_4} = 0.2$ $w_{I_5} = 0.07$ $w_{I_6} = 0.1165$ $w_{I_7} = 0.4$ $P(query) = 1 - P(\neg query)$ $w_{I_7} = 0.4$ $w_{I_7} = 0.4$ $w_{I_7} = 0.4$	$P(query)$ $P(query) = 1 - P(\neg query)$ $P(query)$ $w_{I_1} = 0.0035$ $w_{I_2} = 0.06$ $w_{I_2} = 0.4$ $w_{I_7} = 0.4$ $w_{I_3} = 0.15$ $w_{I_4} = 0.2$ $w_{I_3} = 0.15$ $w_{I_6} = 0.1165$ $w_{I_5} = 0.07$ $w_{I_5} = 0.07$ $w_{I_2} = 0.06$ $w_{I_2} = 0.06$ $w_{I_7} = 0.4$ $w_{I_1} = 0.0035$

	P(query)	$P(query) = 1 - P(\neg query)$		P(query)	$P(query) = 1 - P(\neg query)$
$w_{I_1} = 0.0035$			$w_{I_7}=0.4$		
$w_{I_2}=0.06$			$w_{I_4}=0.2$		
$w_{I_3}=0.15$	0.0635	1 - 0.15 = 0.85	$w_{I_3} = 0.15$	0.2	0.45
$w_{I_4}=0.2$			$w_{I_6} = 0.1165$		
$w_{I_5}=0.07$			$w_{I_5}=0.07$		
$w_{I_6} = 0.1165$			$w_{I_2} = 0.06$		
$w_{I_7}=0.4$			$w_{I_1} = 0.0035$		

	P(query)	$P(query) = 1 - P(\neg query)$		P(query)	$P(query) = 1 - P(\neg query)$
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$w_{I_2}=0.06$			$w_{I_4}=0.2$		
$w_{I_3}=0.15$			$w_{I_3} = 0.15$		
$w_{I_4}=0.2$	0.2635	1 - 0.15 = 0.85	$w_{I_6} = 0.1165$	0.2	0.3335
$w_{I_5} = 0.07$			$w_{I_5}=0.07$		
$w_{I_6} = 0.1165$			$w_{I_2} = 0.06$		
$w_{I_7}=0.4$			$w_{I_1} = 0.0035$		

P(query)

 $w_{I_1} = 0.0035$  $w_{I_2} = 0.06$  $w_{I_3} = 0.15$  $w_{I_4}=0.2$  $w_{I_5} = 0.07$  $w_{I_6} = 0.1165$  $w_{I_7}=0.4$ 

$w_{I_7}=0.4$
$w_{I_4}=0.2$
$w_{I_3}=0.15$
$w_{I_6} = 0.1165$
$w_{I_5}=0.07$
$w_{I_2}=0.06$
$w_{I_1} = 0.0035$

P(query) |  $P(query) = 1 - P(\neg query)$ 

0.2635

1 - 0.7365 = 0.2635

 $P(query) = 1 - P(\neg query)$ 

0.2635

0.2635

### Finding Most Likely Interpretations First



Which value to choose first ? Yes, no, or Sometimes ?  $egin{aligned} \mathrm{Smoking} &\in \{\mathrm{yes}, \mathrm{sometimes}, \mathrm{no}\} \ P(\mathrm{Smoking} = \mathrm{yes}) &= 0.1 \ P(\mathrm{Smoking} = \mathrm{sometimes}) &= 0.3 \ P(\mathrm{Smoking} = \mathrm{no}) &= 0.6 \end{aligned}$ 

### Finding Most Likely Interpretations First



Which value to choose first ? Yes, no, or Sometimes ?  $egin{aligned} \mathrm{Smoking} \in \{\mathrm{yes}, \mathrm{sometimes}, \mathrm{no}\} \ P(\mathrm{Smoking} = \mathrm{yes}) &= 0.1 \ P(\mathrm{Smoking} = \mathrm{sometimes}) &= 0.3 \ P(\mathrm{Smoking} = \mathrm{no}) &= 0.6 \end{aligned}$ 

## Limited Discrepancy Search: Assumptions

Hypothesis 1: The value selection heuristic can be trusted (i.e., favours most likely interpretation first)



### Limited Discrepancy Search: Assumptions

Hypothesis 2: If the heuristics is wrong, it is only at a few nodes of the search tree



### Limited Discrepancy Search: assumptions

Hypothesis 2: If the heuristics is wrong, it is only at a few nodes of the search tree











 $arepsilon = \infty$  [0.0, 1.0]

arepsilon=1.828 $\left[0.1, 0.8\right]$ 




## Mix it all together



## Mix it all together



## Fifty shades of Approximate Inference

	Anytime	Guarantees	Lower bound	Upper bound
d4/gpmc/ exactMC/…	X	Exact on completion	X	X
ApproxMC/ WeightMC	X	$(arepsilon,\delta)$		
PartialKC / SampleSAT		Probability on the LB		
Toulbar		(arepsilon,0)		
Schlandals-LDS		(arepsilon,0)		

## Experiments

- **Marginal Probability in Bayesian Networks**, Reliability estimation in prob. graphs
- Small, large and medium networks
- Timeout: 600s
- Difficult query if solving time > 100s
- Q1. Is Schlandals-LDS efficient ?
- Q2. Better handling of difficult queries ?
- Q3. Does LDS allows faster bound convergence ?

**Experiments: Cactus Plot** 



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# Experiments: Bounds Convergence



#### **Experiments: Bounds Convergence**



#### **Experiments: Bounds Convergence**



### Conclusion & further work

- Deterministic upper- and lower- bounds on the weighted model count
- New Anytime (projected) weighted model counter
- LDS to fasten bounds convergence

#### Conclusion & further work

- Deterministic upper- and lower- bounds on the weighted model count
- New Anytime (projected) weighted model counter
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- Change the discrepancy increment
- Use the bounds to guide the search
- Apply LDS to "classical" weighted model counter

#### Experiments: Cactus Plot



#### **Deterministic Bounds: Lower Bound**

$$WMC(F) = \sum_{j=1}^m P(D_i = d_i^j) imes WMC(F|_{D_i = d_i^j})$$

$$LB = \sum_{j=1}^{m'} P(D_i = d_i^j) imes WMC(F|_{D_i = d_i^j}) \ (m' <= m)$$

#### **Deterministic Bounds: Upper Bound**







 $egin{aligned} \mathrm{Smoking} \in \{\mathrm{yes}, \mathrm{sometimes}, \mathrm{no}\} \ P(\mathrm{Smoking} = \mathrm{yes}) &= 0.1 \ P(\mathrm{Smoking} = \mathrm{sometimes}) &= 0.3 \ P(\mathrm{Smoking} = \mathrm{no}) &= 0.6 \end{aligned}$ 



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### Modelling Distributions in Schlandals



## Branching: classical CNF v.s. Schlandals



